

## Creation, Authentication and

## Recovery of Passwords

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## Outline

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- Introduction
- Master Thesis
- Strength of Passwords
- What we can learn from mistakes


## Identification vs. Authentication

- Password - something you know
- Chip Card/E-Mail - something you have
- Biometrics - something you are
- Password + E-Mail is widely accepted as Authentication!
- Better: Combination of all 3


## Please enter your password

- Computer
- E-Mail
- Online Banking

- Mobile/Smart Phones
- Buildings/Rooms
- ATMs



## Alternatives



- Iris Scan
- Fingerprint
- Chip Cards
- Gestures
- Images
- Voice Analysis


## Problems with Passwords

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- Passwords must/should be
- easy to remember,
- sufficiently long and
- unique (do not reuse passwords).
- Login-Systems must
- Creation
- Authentication
- Recovery
- create and verify passwords,
- provide an option to recover a forgotten password
- and store and transmit passwords in a secure way.


## Master Thesis



## Master Thesis

- Intro
- Strength
- Creating Passwords (RNG, PUF, KDF), Recovery
- Storing Passwords
- Websites (Server)
- Browsers (Client)
- Operating Systems
- Chip Cards


## Master Thesis

- Attacks
- Brute Force
- Dictionaries
- Rainbow Tables
- Alternatives
- KeyPass
- Smartphone + Key Derivation Functions
- Chipcards


## Strength of passwords



It would take
for a desktop PC to crack your password

## Entropy

Definition 1 (Entropy). Let $N$ be the size of our alphabet, the amount of different characters we use (e.g. $N=\#\{a, \ldots, z, A, \ldots, Z, 0,1, \ldots, 9\}=62$ ), and $L$ the length in bit of a password we are trying to measure. The Entropy $H$ is given by

$$
\begin{equation*}
H=\log _{2} N^{L}=L \log _{2} N \tag{2.1}
\end{equation*}
$$

- Compression: "How many bit do we need to store a string using a limited alphabet"
- Here: "How many bit do we need to guess"
- Common: Alphanumeric Alphabet with

$$
\mathbf{N}=26+26+10=62 \text { characters }
$$

- ASCII: 95 printable characters (128 total) Length vs Alphabet - ct'ed
- $2^{\wedge} 5=32$
- $2^{\wedge} 6=64$
- $2^{\wedge} 7=128$
- Chinese?




## Entropy

## Length vs Alphabet - ct'ed 2

- Examples:
- $N=62, L=8: \quad H=47.63$ bit
- $N=62, L=12: \quad H=71.45$ bit
- N=84, L=8: $\quad \mathrm{H}=51.13 \mathrm{bit}$
- N=84, L=12: $\quad \mathrm{H}=76.71 \mathrm{bit}$
- N=95, L=8: $\quad \mathrm{H}=52.56 \mathrm{bit}$
- N=95, L=12: $\mathrm{H}=78.84 \mathrm{bit}$

$$
H=\log _{2} N^{L}=L \log _{2} N
$$

- Increasing length = increasing security?


## Password strength

- Second approach by NIST: Measuring the strength of a password with rules:
- First Character: 4 bit
- Characters 2-8: 2 bit per character
- Characters 9-20: 1.5 bit per character
- Above: 1 bit per character
- Upper + Lower: +6 bit
- Dictionary Search: +6 bit
- Increasing length = increasing redundancy!


## We can learn from mistakes



## UNIX Password Generator (1979)

- System supplied "secure" passwords
- L=8 characters
- Lower case letters and digits ( $\mathrm{N}=36$ )
- Entropy: 41.36 bit (112 years)
- PRNG: 2^15 starting values (Entropy: 15 bit)
R. Morris, K. Thompson: Password Security: A Case History (Communications of the ACM, Volume 22, 1979)


## What we learned from mistakes

- Use PRNG with a sufficiently large seed space


# UNIX Password Store /etc/passwd (197x) 

- Username + Password stored in /etc/passwd
- Later: /etc/shadow + one-way-function
- Everybody on the system could read it
- Everything was fine, until...
\$> ftp
open target.com
Login: ano@nymous.org
get /etc/passwd disconnect


## What we learned from mistakes

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- Use PRNG with a sufficiently large seed space
- Use strong(er) one-way functions to store passwords
- NEVER store passwords in plain text
- OS responsible for restricting access to files


## Windows Password Store LMHASH (1998)

- Max. 14 OEM-characters
- Input: $\mathrm{p}^{\prime}=$ uppercase(substring(p,0,14))
- If less than 14 bytes, add null-bytes;
- Split password into two halves $\mathrm{p}^{\prime}=\mathrm{p} 1| | \mathrm{p} 2$
- Calculate HASH: h = h1 || h2

$$
h_{1}=D E S\left(K G S!@ \# \$ \%, p_{1}\right), h_{2}=\operatorname{DES}\left(K G S!@ \# \$ \%, p_{2}\right)
$$

- Result: 16 byte "hash" value


## Windows Password Store

LMHASH (1998) - ct’ed

- Max. 14 OEM-characters Entropy $(p)<83.4$ bit
- Input: $\mathrm{p}^{\prime}=$ uppercase(substring $(\mathrm{p}, 0,14)$ )

Entropy $\left(p^{\prime}\right)<72.4$ bit

- Assuming alphanumeric numbers, we lost 11 bit of entropy, but 72 bit is still a very good result.


## Windows Password Store LMHASH (1998) - ct'ed 2

- Split password into two halves $\mathrm{p}^{\prime}=\mathrm{p} 1$ || p2
- Calculate "hash": h = h1 || h2 $h_{1}=D E S\left(K G S!@ \# \$ \%, p_{1}\right), h_{2}=D E S\left(K G S!@ \# \$ \%, p_{2}\right)$
- Case 1: Length < 8
$h_{2}=$ DES (KGS!@\#\$\%,0x00000000) =
0xAA 0xD3 0xB4 0x35 0xB5 0x14 0x04 0xEE
- Case 2: Length >= 8
$\operatorname{Entropy}\left(p_{1}\right)=\operatorname{Entropy}\left(p_{2}\right) \leqslant 7 \log _{2} 36=36.2$ bit


## Windows Password Store LMHASH (1998) - ct'ed 3

- Split password into two halves $\mathrm{p}^{\prime}=\mathrm{p} 1$ || p2
- Calculate "hash": h = h1 || h2 $h_{1}=D E S\left(K G S!@ \# \$ \%, p_{1}\right), h_{2}=D E S\left(K G S!@ \# \$ \%, p_{2}\right)$
- Case 2: Length >= 8
$\operatorname{Entropy}\left(p_{1}\right)=\operatorname{Entropy}\left(p_{2}\right) \leqslant 7 \log _{2} 36=36.2$ bit
- Instead of $\log _{2}\left(N^{14}\right)$ we now have $\log _{2}\left(2 \cdot N^{7}\right)=1+\log _{2}\left(N^{7}\right)=1+7 \log _{2} 36=37.2$ bit


## What we learned from mistakes

- Use PRNG with a sufficiently large seed space
- Use strong(er) one-way functions to store passwords
- NEVER store passwords in plain text
- OS responsible for restricting access to files
- Microsoft (20 years later): Use strong(er) oneway functions for authentication


## EuroCheque ATM PINs 1981-1997 (Germany)

## PIN Calculation for EuroCheque ATM Debit Cards



- 1997: M. Kuhn: Probability Theory for Pickpockets - ec-PIN guessing
- Showed that success probability for breaking in can be increased from $0.03 \%$ to $0.7 \%$


## More mistakes...



## Input Devices

- Smudge Attack
- Thermal Imaging


Most used passwords

| $\#$ | Password | $\#$ | Password |  | $\#$ | PIN |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | password | 14 | sunshine |  | 1234 |  |
| 2 | 123456 | 15 | master |  | 12 | 0000 |
| 3 | 12345678 | 16 | 123123 |  | 3 | 2580 |
| 4 | abc123 | 17 | welcome |  | 4 | 1111 |
| 5 | qwerty | 18 | shadow |  | 5 | 5555 |
| 6 | monkey | 19 | ashley |  | 6 | 5683 |
| 7 | letmein | 20 | football |  | 7 | 0852 |
| 8 | dragon | 21 | jesus |  | 8 | 2222 |
| 9 | 111111 | 22 | michael |  | 9 | 1212 |
| 10 | baseball | 23 | ninja |  | 10 | 1998 |
| 11 | iloveyou | 24 | mustang |  |  |  |
| 12 | trustno1 | 25 | password1 |  |  |  |
| 13 | 1234567 |  |  |  |  |  |

## Most used PINs



| $\#$ | PIN |
| :--- | :--- |
| 1 | 1234 |
| 2 | 0000 |
| 3 | 2580 |
| 4 | 1111 |
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